

## $\mathcal{V}$ -CAT IS LOCALLY PRESENTABLE OR LOCALLY BOUNDED IF $\mathcal{V}$ IS SO

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ABSTRACT. We show, for a monoidal closed category  $\mathcal{V} = (\mathcal{V}_0, \otimes, I)$ , that the category  $\mathcal{V}\text{-Cat}$  of small  $\mathcal{V}$ -categories is locally  $\lambda$ -presentable if  $\mathcal{V}_0$  is so, and that it is locally  $\lambda$ -bounded if the closed category  $\mathcal{V}$  is so, meaning that  $\mathcal{V}_0$  is locally  $\lambda$ -bounded and that a side condition involving the monoidal structure is satisfied.

Many important properties of a monoidal category  $\mathcal{V}$  are inherited by the category  $\mathcal{V}\text{-Cat}$  of small  $\mathcal{V}$ -categories. For instance, if  $\mathcal{V}$  is symmetric monoidal,  $\mathcal{V}\text{-Cat}$  has a canonical symmetric monoidal structure, as was observed already in [4]. Much later [7, Remark 5.2], it was realized that if  $\mathcal{V}$  is only *braided* monoidal then  $\mathcal{V}\text{-Cat}$  still has a canonical monoidal structure, although it need not have a braiding unless the braiding on  $\mathcal{V}$  is in fact a symmetry. Similarly, it is straightforward to show that  $\mathcal{V}\text{-Cat}$  is monoidal *closed* when  $\mathcal{V}$  is closed and complete, and that  $\mathcal{V}\text{-Cat}$  is complete when  $\mathcal{V}$  is so. All of these results are essentially routine; the less trivial fact that  $\mathcal{V}\text{-Cat}$  is cocomplete when  $\mathcal{V}$  is so was first proved in [11].

The properties of  $\mathcal{V}$  or  $\mathcal{V}\text{-Cat}$  that we consider here are of a less basic nature, being conditions on  $\mathcal{V}$  which allow proofs by transfinite induction of the existence of various important adjoints. The best known of these conditions is *local presentability* [5], but there is also the notion of *local boundedness* [8], which is more general than local presentability, but also much more common, and sufficient for the central existence results of [8, Chapter 6], from which follow the basic results of the theory of enriched projective sketches. Recall that to be locally presentable is to be locally  $\lambda$ -presentable for some regular cardinal  $\lambda$ , and similarly that to be locally bounded is to be locally  $\lambda$ -bounded for some  $\lambda$ . It would be one thing to prove that  $\mathcal{V}\text{-Cat}$  is locally presentable if  $\mathcal{V}$  is so (in the sense that its underlying ordinary category  $\mathcal{V}_0$  is so); here we prove the stronger result that  $\mathcal{V}\text{-Cat}$  is locally  $\lambda$ -presentable if  $\mathcal{V}_0$  is so, so that the passage from  $\mathcal{V}$  to  $\mathcal{V}\text{-Cat}$  does not require the regular cardinal  $\lambda$  to be changed. When it comes to local boundedness, we prove that  $\mathcal{V}\text{-Cat}$  is locally  $\lambda$ -bounded when  $\mathcal{V}$  is so “as a closed category”, meaning that  $\mathcal{V}_0$  is locally  $\lambda$ -bounded and satisfies a side condition involving the monoidal structure. We recall the precise definitions of local  $\lambda$ -presentability and local  $\lambda$ -boundedness in Section 2, but the common aspect is that  $\mathcal{V}_0$  is cocomplete and has a small set  $\mathcal{G}$  of objects forming in some sense a generator of  $\mathcal{V}_0$ , with the representables  $\mathcal{V}_0(G, -) : \mathcal{V}_0 \rightarrow \mathbf{Set}$  preserving certain colimits:  $\lambda$ -filtered colimits in the locally  $\lambda$ -presentable case, and  $\lambda$ -filtered unions

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(with respect to a given factorization system on  $\mathcal{V}_0$ ) in the locally  $\lambda$ -bounded case.

All categories are assumed to have small hom-sets.

### 1. Unions

For this section we consider a cocomplete category  $\mathcal{K}$  with a proper factorization system  $(\mathcal{E}, \mathcal{M})$ ; recall that  $(\mathcal{E}, \mathcal{M})$  is *proper* when each  $\mathcal{E}$  is an epimorphism and each  $\mathcal{M}$  a monomorphism — equivalently, when each  $\mathcal{M}$  is a monomorphism and each coretraction is in  $\mathcal{M}$ . Note that a map  $f : A \rightarrow B$  lies in  $\mathcal{E}$  precisely when each factorization  $f = mg$  with  $m \in \mathcal{M}$  has  $m$  invertible.

A small family  $(m_j : A_j \rightarrow B)_{j \in J}$  of maps with a common codomain is said to be *jointly in  $\mathcal{E}$*  if the induced map  $m : \sum_j A_j \rightarrow B$  is in  $\mathcal{E}$ ; this is equivalent to saying that there is no proper  $\mathcal{M}$ -subobject of  $B$  through which each  $m_j$  factorizes. When moreover each  $m_j$  is in  $\mathcal{M}$ , we say that the family constitutes an  *$\mathcal{M}$ -union*, or that  $B$  is the  $\mathcal{M}$ -union of the  $m_j$ .

More generally, the  $\mathcal{M}$ -union of a small family  $(m_j : A_j \rightarrow B)_{j \in J}$  of  $\mathcal{M}$ -subobjects of  $B$  is defined to be the unique  $\mathcal{M}$ -subobject  $n : A \rightarrow B$  containing the  $m_j$  for which the corresponding  $n_j : A_j \rightarrow A$  constitute an  $\mathcal{M}$ -union. We may calculate this  $\mathcal{M}$ -union by taking the  $(\mathcal{E}, \mathcal{M})$ -factorization of  $m : \sum_j A_j \rightarrow B$ .

We shall need to speak of *preservation of  $\mathcal{M}$ -unions* only in the case of representable functors. We say that the representable functor  $\mathcal{K}(X, -) : \mathcal{K} \rightarrow \mathbf{Set}$  preserves the  $\mathcal{M}$ -union  $(m_j : A_j \rightarrow B)_{j \in J}$  if the functions  $\mathcal{K}(X, m_j) : \mathcal{K}(X, A_j) \rightarrow \mathcal{K}(X, B)$  are jointly surjective; in more concrete terms this says that any map  $f : X \rightarrow B$  factorizes through some  $m_j$ .

Given a small family  $(m_j : A_j \rightarrow B)_{j \in J}$  we can preorder the set  $J$  by setting  $j \leq k$  whenever  $A_j \leq A_k$  as  $\mathcal{M}$ -subobjects of  $B$ . Then the  $A_j$  are the object values of a functor  $A : J \rightarrow \mathcal{K}$ , and we may form  $\text{colim} A$  and the induced  $h : \text{colim} A \rightarrow B$ . It is easy to see that the  $m_j$  are an  $\mathcal{M}$ -union if and only if  $h \in \mathcal{E}$ .

Finally for a regular cardinal  $\lambda$ , the preorder  $J$  is said to be  $\lambda$ -filtered if it is so as a category: that is, if for each subset  $K$  of  $J$  with cardinality less than  $\lambda$ , the  $A_k$  with  $k \in K$  are all contained in some  $A_j$ . By a  $\lambda$ -filtered  $\mathcal{M}$ -union  $(m_j : A_j \rightarrow B)_{j \in J}$  we mean one for which  $J$  is  $\lambda$ -filtered.

### 2. Locally presentable and locally bounded categories

In this section we continue to consider a cocomplete category  $\mathcal{K}$ ; from time to time we shall further suppose it to be equipped with a proper factorization system  $(\mathcal{E}, \mathcal{M})$ .

Let  $\lambda$  be a regular cardinal. An object  $X$  of  $\mathcal{K}$  is said to be  $\lambda$ -presentable [5] if the representable functor  $\mathcal{K}(X, -) : \mathcal{K} \rightarrow \mathbf{Set}$  preserves  $\lambda$ -filtered colimits, and  $\lambda$ -bounded [6] if  $\mathcal{K}(X, -)$  preserves  $\lambda$ -filtered  $\mathcal{M}$ -unions.

A small set  $\mathcal{G}$  of objects of  $\mathcal{K}$  is said to be a *strong generator* if an arrow  $f : A \rightarrow B$  is invertible whenever  $\mathcal{K}(G, f) : \mathcal{K}(G, A) \rightarrow \mathcal{K}(G, B)$  is bijective for each  $G \in \mathcal{G}$ ; while

$\mathcal{G}$  is an  $(\mathcal{E}, \mathcal{M})$ -generator if this is true for arrows  $f : A \rightarrow B$  in  $\mathcal{M}$ . Clearly  $\mathcal{G}$  is an  $(\mathcal{E}, \mathcal{M})$ -generator precisely when, for each  $A \in \mathcal{K}$ , the family of all maps  $G \rightarrow A$  with  $G \in \mathcal{G}$  is jointly in  $\mathcal{E}$ ; that is, when the evident map  $\epsilon_A : \sum_{G \in \mathcal{G}} \mathcal{K}(G, A) \bullet G \rightarrow A$  lies in  $\mathcal{E}$ ; here we are writing  $X \bullet A$  for the coproduct of  $X$  copies of  $A$ . When  $(\mathcal{E}, \mathcal{M})$  is the proper factorization system (strong epimorphisms, monomorphisms), we have the well-known result (see for instance [6, Proposition 2.5.3]) that an  $(\mathcal{E}, \mathcal{M})$ -generator is the same thing as a strong generator. In our cocomplete category  $\mathcal{K}$ , the pair (strong epimorphisms, monomorphisms) is certainly a proper factorization system if  $\mathcal{K}$  admits arbitrary cointersections of strong epimorphisms.

The cocomplete  $\mathcal{K}$  is said to be *locally  $\lambda$ -presentable* if it has a strong generator all of whose objects are  $\lambda$ -presentable; it is a consequence that  $\mathcal{K}$  is then complete. This and many other facts about locally presentable categories can be found in the books [1, 5, 10].

The cocomplete  $\mathcal{K}$  is said to be *locally  $\lambda$ -bounded* with respect to a proper factorization system  $(\mathcal{E}, \mathcal{M})$  if it has an  $(\mathcal{E}, \mathcal{M})$ -generator all of whose objects are  $\lambda$ -bounded, and if moreover  $\mathcal{K}$  admits arbitrary cointersections (even large ones, if need be) of maps in  $\mathcal{E}$ . The definition of locally  $\lambda$ -bounded category given in [8] included the further assumption of completeness, but once again this is a consequence of the other axioms, as we show in Corollary 2.2 below.

As well as being complete, every locally  $\lambda$ -presentable category is well-powered; it follows that it has a proper factorization system  $(\mathcal{E}, \mathcal{M})$  in which  $\mathcal{M}$  consists of the monomorphisms and  $\mathcal{E}$  the strong epimorphisms. For this factorization system, an  $(\mathcal{E}, \mathcal{M})$ -generator is, as we observed above, the same thing as a strong generator. Locally presentable categories are also well-copowered, and so arbitrary  $\mathcal{E}$ -cointersections exist. Finally, it turns out (see [6, Lemma 2.3.1]) that in a locally  $\lambda$ -presentable category every  $\lambda$ -presentable object is  $\lambda$ -bounded; we deduce that every locally  $\lambda$ -presentable category is locally  $\lambda$ -bounded. The converse, however, is false: see [5, p.104] or [6, p.190] for examples of locally  $\lambda$ -bounded categories that are not locally  $\mu$ -presentable for any  $\mu$ .

A cocomplete monoidal closed category is said to be *locally  $\lambda$ -bounded as a closed category* if its underlying ordinary category is locally  $\lambda$ -bounded and, in addition, the functors  $A \otimes -$  and  $- \otimes A$  map  $\mathcal{E}$  into  $\mathcal{E}$  for all objects  $A$ . The latter condition is clearly equivalent to the condition that  $e \otimes e' \in \mathcal{E}$  whenever  $e, e' \in \mathcal{E}$ , and it turns out to be vacuous if  $\mathcal{M}$  consists of all the monomorphisms.

In fact all the examples of closed categories considered in [8] have some factorization system for which they are locally bounded. Algebraic examples, such as the categories **Set**, **Cat**, and **Ab** of sets, categories, and abelian groups are all locally finitely presentable, as is the combinatorial example **SSet**, the category of simplicial sets. The reason for using the weaker notion of local boundedness rather than local presentability is the desire to include such topological examples as the categories **CGTop**, **QTop**, and **Ban** of compactly generated topological spaces, quasi-topological spaces, and Banach spaces, which are not locally presentable, but are locally bounded. The example **QTop** is not  $\mathcal{E}$ -wellcopowered, which explains why we must explicitly require arbitrary cointersections of maps in  $\mathcal{E}$ . For the details, and for many further examples, including Lawvere's closed category given by the interval  $[0, \infty]$  of the reals, see [8, Chapter 6].

For our promised proof that every locally bounded category is complete we use an (apparently unpublished)  $(\mathcal{E}, \mathcal{M})$ -variant of Freyd’s Special Adjoint Functor Theorem, namely:

**2.1. PROPOSITION.** *Let the cocomplete category  $\mathcal{K}$  have the factorization system  $(\mathcal{E}, \mathcal{M})$  for which  $\mathcal{E}$  is contained in the epimorphisms; suppose that  $\mathcal{K}$  admits arbitrary cointersections of maps in  $\mathcal{E}$ , and that  $\mathcal{K}$  has an  $(\mathcal{E}, \mathcal{M})$ -generator  $\mathcal{G}$ . Then every cocontinuous functor  $S : \mathcal{K} \rightarrow \mathcal{L}$  has a right adjoint.*

**PROOF.** To provide a right adjoint to  $S$  is equally to provide, for each  $D \in \mathcal{L}$ , a terminal object of the comma category  $S/D$ , whose objects are pairs  $(C, f : SC \rightarrow D)$  and whose maps  $(C, f) \rightarrow (C', f')$  are maps  $x : C \rightarrow C'$  with  $Sx.f = f'$ . The forgetful functor  $U : S/D \rightarrow \mathcal{K}$  creates colimits (and hence reflects epimorphisms). We get an induced factorization system, still called  $(\mathcal{E}, \mathcal{M})$ , on  $S/D$  by taking  $x : (C, f) \rightarrow (C', f')$  to be in  $\mathcal{E}$  or in  $\mathcal{M}$  when  $Ux$  is so; once again every  $\mathcal{E}$  is an epimorphism. Finally, the small set consisting of the  $(C, f)$  with  $C \in \mathcal{G}$  forms an  $(\mathcal{E}, \mathcal{M})$ -generator for  $S/D$ . Thus  $(S/D, \mathcal{E}, \mathcal{M})$  has just the properties required in the proposition of  $(\mathcal{K}, \mathcal{E}, \mathcal{M})$ . So it suffices to prove that the  $\mathcal{K}$  of the proposition has a terminal object.

Form in  $\mathcal{K}$  the coproduct  $H = \sum_{G \in \mathcal{G}} G$ , and let  $\zeta : H \rightarrow K$  be the cointersection of all the maps in  $\mathcal{E}$  having domain  $H$ ; of course  $\zeta \in \mathcal{E}$  and is an epimorphism. Any two maps  $f, g : A \rightarrow K$  must coincide: for their coequalizer  $h : K \rightarrow L$  is in  $\mathcal{E}$ , so that  $h\zeta$  is in  $\mathcal{E}$ , whence  $kh\zeta = \zeta$  for some  $k$  by the definition of  $\zeta$  as the smallest  $\mathcal{E}$ -quotient, so that in fact  $kh = 1$  and  $h$  is invertible.

To exhibit  $K$  as the desired terminal object it remains only to show that, for each  $A \in \mathcal{K}$ , there is a map  $A \rightarrow K$ . For each  $G \in \mathcal{G}$  and  $A \in \mathcal{K}$  we have the trivial function  $\mathcal{K}(G, A) \rightarrow 1$  into the singleton set, so that we have an induced map  $t : \sum_{G \in \mathcal{G}} \mathcal{K}(G, A) \bullet G \rightarrow \sum_{G \in \mathcal{G}} G$ . Form in  $\mathcal{C}$  the pushout

$$\begin{array}{ccc} \sum_{G \in \mathcal{G}} \mathcal{K}(G, A) \bullet G & \xrightarrow{\epsilon_A} & A \\ \downarrow t & & \downarrow r \\ \sum_{G \in \mathcal{G}} G & \xrightarrow{s} & L ; \end{array}$$

here  $\epsilon_A$  lies in  $\mathcal{E}$  since  $\mathcal{G}$  is an  $(\mathcal{E}, \mathcal{M})$ -generator, so that its pushout  $s$  also lies in  $\mathcal{E}$ . By the definition of  $K$ , therefore, there is a map  $v : L \rightarrow K$ , and thus a map  $vr : A \rightarrow K$ . ■

**2.2. COROLLARY.** *Let the cocomplete category  $\mathcal{K}$  have a factorization system  $(\mathcal{E}, \mathcal{M})$  for which every  $\mathcal{E}$  is an epimorphism, and suppose that  $\mathcal{K}$  admits arbitrary cointersections of maps in  $\mathcal{E}$  and has an  $(\mathcal{E}, \mathcal{M})$ -generator  $\mathcal{G}$ . Then  $\mathcal{K}$  is complete.*

**PROOF.** For each small category  $\mathcal{C}$  we seek a right adjoint to the diagonal  $\Delta : \mathcal{K} \rightarrow [\mathcal{C}, \mathcal{K}]$ ; and this adjoint exists by the proposition, since  $[\mathcal{C}, \mathcal{K}]$  has colimits formed pointwise and  $\Delta$  is cocontinuous. ■

2.3. REMARK. Given a cocomplete category  $\mathcal{K}$ , to give a factorization system  $(\mathcal{E}, \mathcal{M})$  having each  $\mathcal{E}$  epimorphic and admitting arbitrary cointersections of maps in  $\mathcal{E}$ , it suffices by [3, Lemma 3.1] to give a class  $\mathcal{E}$  of epimorphisms in  $\mathcal{K}$ , closed under composition and stable under pushout, for which arbitrary cointersections of maps in  $\mathcal{E}$  exist and lie in  $\mathcal{E}$ .

Before leaving this section, we make a final observation of rather lesser importance. We have discussed what it means for a monoidal closed category to be locally bounded *as a closed category*, but we have not considered local presentability for closed categories. In [9], a monoidal closed category  $\mathcal{V}$  was defined to be *locally  $\lambda$ -presentable as a closed category* if its underlying category  $\mathcal{V}_0$  was locally  $\lambda$ -presentable and the  $\lambda$ -presentable objects of  $\mathcal{V}_0$  were closed under the monoidal structure: that is, the unit  $I$  was  $\lambda$ -presentable and  $X \otimes Y$  was  $\lambda$ -presentable whenever  $X$  and  $Y$  were so. The observation we wish to make here is the following:

2.4. PROPOSITION. *If  $\mathcal{V}$  is a monoidal closed category and  $\mathcal{V}_0$  is locally  $\lambda$ -presentable, then there exists a regular cardinal  $\mu$  for which  $\mathcal{V}$  is locally  $\mu$ -presentable as a closed category.*

PROOF. Observe that the set of  $\lambda$ -presentable objects is (essentially) small, so the set of objects of the form  $G \otimes H$  where  $G$  and  $H$  are  $\lambda$ -presentable is (essentially) small. Thus there exists a regular cardinal  $\mu$  with the property that  $I$  is  $\mu$ -presentable and  $G \otimes H$  is  $\mu$ -presentable whenever  $G$  and  $H$  are  $\lambda$ -presentable. But now if  $A$  and  $B$  are  $\mu$ -presentable objects, then we may write  $A = \text{colim}_i G_i$  and  $B = \text{colim}_j H_j$  where the colimits in question are  $\mu$ -small, and where each  $G_i$  and each  $H_j$  is  $\lambda$ -presentable. Then

$$\begin{aligned} A \otimes B &= \text{colim}_i G_i \otimes \text{colim}_j H_j \\ &= \text{colim}_{i,j} (G_i \otimes H_j) \end{aligned}$$

and each  $G_i \otimes H_j$  is  $\mu$ -presentable; thus  $A \otimes B$  is a  $\mu$ -small colimit of  $\mu$ -presentable objects, and thus is itself  $\mu$ -presentable. This proves that  $\mathcal{V}$  is locally  $\mu$ -presentable as a closed category. ■

### 3. $\mathcal{V}$ -Cat is finitarily monadic over $\mathcal{V}$ -Gph

For this section we suppose that  $\mathcal{V}$  is a monoidal category which is cocomplete, and that the functors  $A \otimes -$  and  $- \otimes A$  preserve colimits for all objects  $A$  of  $\mathcal{V}$ , as is certainly the case if the monoidal  $\mathcal{V}$  is *closed*.

As a preliminary to our investigation of  $\mathcal{V}$ -Cat, we consider the category  $\mathcal{V}$ -Gph of  $\mathcal{V}$ -graphs and their morphisms. Recall that a  $\mathcal{V}$ -graph is a pair  $(X, A)$ , where  $X$  is a (small) set, and  $A$  is a family  $(A(x, y))_{x,y \in X}$  of objects of  $\mathcal{V}$ . A  $\mathcal{V}$ -graph morphism from  $(X, A)$  to  $(Y, B)$  is a pair  $(f, \varphi)$  where  $f : X \rightarrow Y$  is a function from  $X$  to  $Y$ , and  $\varphi$  is a family  $(\varphi_{x,y} : A(x, y) \rightarrow B(fx, fy))_{x,y \in X}$  of morphisms in  $\mathcal{V}$ . We write  $P : \mathcal{V}\text{-Gph} \rightarrow \mathbf{Set}$  for the functor sending a  $\mathcal{V}$ -graph  $(X, A)$  to its set  $X$  of objects, and sending  $(f, \varphi)$  to  $f$ .

There is an evident forgetful functor  $U : \mathcal{V}\text{-Cat} \rightarrow \mathcal{V}\text{-Gph}$  which is monadic, as was proved in [2] under the hypotheses above, and more generally when  $\mathcal{V}$  is a suitable

bicategory; and much earlier in [11] when  $\mathcal{V}$  is symmetric monoidal closed. In this section we shall show that the monad in question is finitary — meaning that it preserves filtered colimits; in the next, we show that  $\mathcal{V}\text{-Gph}$  is locally  $\lambda$ -presentable if  $\mathcal{V}$  is so; it will then follow that  $\mathcal{V}\text{-Cat}$  is locally  $\lambda$ -presentable if  $\mathcal{V}$  is so, by [5, Satz 10.3]. Accordingly we begin by studying colimits in  $\mathcal{V}\text{-Gph}$ .

Following [2], we shall analyze  $\mathcal{V}$ -graphs in terms of the more general  $\mathcal{V}$ -matrices. If  $X$  and  $Y$  are sets, a  $\mathcal{V}$ -matrix  $S$  from  $X$  to  $Y$  is a family  $(S(y, x))_{(x,y) \in X \times Y}$  of objects of  $\mathcal{V}$ ; thus a  $\mathcal{V}$ -graph is just a set  $X$  equipped with a  $\mathcal{V}$ -matrix  $A : X \rightarrow X$ . The value of  $\mathcal{V}$ -matrices is that they can be composed: if  $S : X \rightarrow Y$  and  $T : Y \rightarrow Z$  are  $\mathcal{V}$ -matrices, then their composite  $TS : X \rightarrow Z$  is defined by

$$(TS)(z, x) = \sum_{y \in Y} T(z, y) \otimes S(y, x).$$

There is now a bicategory  $\mathcal{V}\text{-Mat}$  in which the objects are the (small) sets, the 1-cells are the  $\mathcal{V}$ -matrices, and a 2-cell between  $\mathcal{V}$ -matrices  $S, S' : X \rightarrow Y$  is a family  $(\sigma_{y,x} : S(y, x) \rightarrow S'(y, x))_{(x,y) \in X \times Y}$  of morphisms of  $\mathcal{V}$ .

For objects  $X$  and  $Y$  of  $\mathcal{V}\text{-Mat}$ , the hom-category  $\mathcal{V}\text{-Mat}(X, Y)$  is just  $\mathcal{V}^{Y \times X}$ , which is cocomplete since  $\mathcal{V}$  is so, with colimits formed pointwise from those in  $\mathcal{V}$ . Furthermore, if  $S : Y \rightarrow Y'$  and  $R : X' \rightarrow X$  are arbitrary  $\mathcal{V}$ -matrices, the functors  $\mathcal{V}\text{-Mat}(X, S) : \mathcal{V}\text{-Mat}(X, Y) \rightarrow \mathcal{V}\text{-Mat}(X, Y')$  and  $\mathcal{V}\text{-Mat}(R, Y) : \mathcal{V}\text{-Mat}(X, Y) \rightarrow \mathcal{V}\text{-Mat}(X', Y)$  are cocontinuous; we express this fact by saying that “composition commutes with colimits”.

A function  $f : X \rightarrow Y$  determines  $\mathcal{V}$ -matrices  $f_* : X \rightarrow Y$  and  $f^* : Y \rightarrow X$  with

$$f_*(y, x) = f^*(x, y) = \begin{cases} I & \text{if } fx = y \\ 0 & \text{otherwise} \end{cases}$$

where  $I$  denotes the unit object and  $0$  the initial object of  $\mathcal{V}$ . The reader will easily construct a natural bijection between 2-cells  $f_*A \rightarrow B$  and 2-cells  $A \rightarrow f^*B$ , and so deduce that  $f_*$  is left adjoint to  $f^*$  in the bicategory  $\mathcal{V}\text{-Mat}$ . In fact it is also easy to describe explicitly the unit  $1_X \rightarrow f^*f_*$  and the counit  $f_*f^* \rightarrow 1_Y$ .

We have already observed that a  $\mathcal{V}$ -graph is an object  $X$  of  $\mathcal{V}\text{-Mat}$  equipped with a 1-cell  $A : X \rightarrow X$ ; a morphism of  $\mathcal{V}$ -graphs from  $(X, A)$  to  $(Y, B)$  can be seen as a function  $f : X \rightarrow Y$  equipped with a 2-cell  $\varphi : A \rightarrow f^*Bf_*$ , as the following calculation shows:

$$\begin{aligned} (f^*Bf_*)(z, x) &= \sum_{y \in Y} f^*(z, y) \otimes (Bf_*)(y, x) \\ &= (Bf_*)(fz, x) \\ &= \sum_{y \in Y} B(fz, y) \otimes f_*(y, x) \\ &= B(fz, fx). \end{aligned}$$

In fact, because of the adjunction  $f_* \dashv f^*$  in the bicategory  $\mathcal{V}\text{-Mat}$ , there is a bijection (of “mates”) between 2-cells  $\varphi : A \rightarrow f^*Bf_*$  and 2-cells  $\widehat{\varphi} : f_*Af^* \rightarrow B$ ; explicitly, we find that

$$(f_*Af^*)(u, v) = \sum_{\substack{fx=u \\ fy=v}} A(x, y),$$

and now for  $x \in f^{-1}(u)$  and  $y \in f^{-1}(v)$  the  $(x, y)$ -component of  $\widehat{\varphi}_{u,v} : (f_*Af^*)(u, v) \rightarrow B(u, v)$  is  $\varphi_{x,y}$ .

As shown in [2], colimits in  $\mathcal{V}\text{-Gph}$  can be described as follows. Let  $\mathcal{J}$  be a small category, and  $(X, A) : \mathcal{J} \rightarrow \mathcal{V}\text{-Gph}$  a functor; we denote the image of an object  $j$  under  $(X, A)$  by  $(X_j, A_j)$  and the image of a morphism  $\theta : j \rightarrow k$  by  $(X_\theta, A_\theta)$ . Consider the functor  $X = P(X, A) : \mathcal{J} \rightarrow \mathbf{Set}$ , and form its colimit  $\bar{X}$  with colimit cone  $(q_j : X_j \rightarrow \bar{X})_{j \in \mathcal{J}}$ . There is a functor  $\tilde{A} : \mathcal{J} \rightarrow \mathcal{V}\text{-Mat}(\bar{X}, \bar{X})$  sending  $j$  to  $(q_j)_*A_j(q_j)^*$  and sending a morphism  $\theta : j \rightarrow k$  to  $(q_k)_*\widehat{A}_\theta(q_k)^* : (q_k)_*(X_\theta)_*A_j(X_\theta)^*(q_k)^* \rightarrow (q_k)_*A_k(q_k)^*$ , where  $\widehat{A}_\theta : (X_\theta)_*A_j(X_\theta)^* \rightarrow A_k$  is the mate, as above, of  $A_\theta : A_j \rightarrow (X_\theta)^*A_k(X_\theta)_*$ . As we saw above, the colimit of  $\tilde{A}$  is formed pointwise from colimits in  $\mathcal{V}$ : write  $\bar{A} : \bar{X} \rightarrow \bar{X}$  for this colimit, with colimit cone  $\alpha'_j : (q_j)_*A_j(q_j)^* \rightarrow \bar{A}$ . Now we have in  $\mathcal{V}\text{-Mat}$  a cone  $(q_j, \alpha_j) : (X_j, A_j) \rightarrow (\bar{X}, \bar{A})$ , where  $\alpha_j : A_j \rightarrow (q_j)^*\bar{A}(q_j)_*$  is the 2-cell for which  $\widehat{\alpha}_j : (q_j)_*A_j(q_j)^* \rightarrow \bar{A}$  is  $\alpha'_j$ ; and it is shown in [2] that this is a colimit cone for  $(X, A) : \mathcal{J} \rightarrow \mathcal{V}\text{-Gph}$ . (Of course we henceforth drop the name  $\alpha'_j$  in favour of  $\widehat{\alpha}_j$ .)

We need below to consider functors  $(X, A) : \mathcal{J} \rightarrow \mathcal{V}\text{-Gph}$  and  $(X, B) : \mathcal{J} \rightarrow \mathcal{V}\text{-Gph}$  with the same  $X : \mathcal{J} \rightarrow \mathbf{Set}$ ; accordingly we introduce the category  $\mathcal{V}\text{-Gph}^{(2)}$  defined by the pullback

$$\begin{array}{ccc} \mathcal{V}\text{-Gph}^{(2)} & \xrightarrow{Q} & \mathcal{V}\text{-Gph} \\ R \downarrow & & \downarrow P \\ \mathcal{V}\text{-Gph} & \xrightarrow{P} & \mathbf{Set} \end{array}$$

in  $\mathbf{Cat}$ ; observe that, since  $\mathcal{V}\text{-Gph}$  and  $\mathbf{Set}$  are cocomplete and  $P$  is cocontinuous,  $\mathcal{V}\text{-Gph}^{(2)}$  is cocomplete and the functors  $Q$  and  $R$  jointly create colimits. An object of  $\mathcal{V}\text{-Gph}^{(2)}$  is a pair  $((X, A), (X, B))$  of  $\mathcal{V}$ -graphs with the same underlying set  $X$ , which we henceforth write as  $(X, A, B)$ ; and a morphism has the form  $(f, \alpha, \beta) : (X, A, B) \rightarrow (X', A', B')$  where  $(f, \alpha) : (X, A) \rightarrow (X', A')$  and  $(f, \beta) : (X, B) \rightarrow (X', B')$  are morphisms in  $\mathcal{V}\text{-Gph}$ . To give a pair of functors as in the first sentence of this paragraph is of course to give a single functor from  $\mathcal{J}$  to  $\mathcal{V}\text{-Gph}^{(2)}$ . In the same way we can define  $\mathcal{V}\text{-Gph}^{(n)}$  with objects  $(X, A_1, \dots, A_n)$  by taking the fibred product in  $\mathbf{Cat}$  of  $n$  copies of  $P : \mathcal{V}\text{-Gph} \rightarrow \mathbf{Set}$ , and  $\mathcal{V}\text{-Gph}^{(\mathbb{N})}$  by taking the fibred product of copies indexed by the set  $\mathbb{N}$  of natural numbers; and we have the corresponding results about colimits in  $\mathcal{V}\text{-Gph}^{(n)}$  and  $\mathcal{V}\text{-Gph}^{(\mathbb{N})}$ .

Consider the functor  $S : \mathcal{V}\text{-Gph}^{(2)} \rightarrow \mathcal{V}\text{-Gph}$  sending  $(X, A, B)$  to  $(X, A + B)$ , where the sum  $A + B$  of matrices is of course the coproduct in  $\mathcal{V}^{X \times X}$ ; the value of  $S$  on morphisms is given by the evident sum of 2-cells using the distributive law for matrices.

This functor  $S$  preserves colimits, for if  $(X, A, B) : \mathcal{J} \rightarrow \mathcal{V}\text{-Gph}^{(2)}$ , it is clear from the description above of colimits in  $\mathcal{V}\text{-Gph}$  that the colimit of  $(X, A + B)$  is  $(\bar{X}, \bar{A} + \bar{B})$ , where  $(\bar{X}, \bar{A})$  and  $(\bar{X}, \bar{B})$  are the colimits of  $(X, A)$  and  $(X, B)$ . Similarly of course for sums of any size: the form we need below is:

**3.1. LEMMA.** *The functor  $S : \mathcal{V}\text{-Gph}^{(\mathbb{N})} \rightarrow \mathcal{V}\text{-Gph}$  sending  $(X, (A_n)_{n \in \mathbb{N}})$  to  $(X, \sum_{n \in \mathbb{N}} A_n)$  preserves colimits. ■*

We also need to consider the functor  $M : \mathcal{V}\text{-Gph}^{(2)} \rightarrow \mathcal{V}\text{-Gph}$  which sends  $(X, A, B)$  to  $(X, AB)$ , where  $AB$  denotes as before the matrix product. We must of course define  $M$  on morphisms too. Recall that the  $\alpha$  of a morphism  $(f, \alpha) : (X, A) \rightarrow (X', A')$  can be seen as a matrix  $\alpha : A \rightarrow f^*A'f_*$ , but can equally be described by its mate  $\hat{\alpha} : f_*Af^* \rightarrow A'$  under the adjunction  $f_* \dashv f^*$ . But there is of course yet another equivalent form, namely  $\bar{\alpha} : f_*A \rightarrow A'f_*$ . In fact we find that  $(f_*A)(x', x) = \sum_{fy=x'} A(y, x)$ , that  $(A'f_*)(x', x) = A'(x', fx)$ , and that  $\bar{\alpha}_{x',x}$  has  $\alpha_{y,x}$  as its  $y$ -component. Now the value of  $M$  on  $(f, \alpha, \beta) : (X, A, B) \rightarrow (X', A', B')$  is  $(f, \gamma) : (X, AB) \rightarrow (X', A'B')$  where  $\gamma$  is determined in terms of its mate  $\bar{\gamma}$  by the pasting composite

$$\begin{array}{ccc}
 X & \xrightarrow{f_*} & X' \\
 \downarrow AB & \cong & \downarrow A'B' \\
 X & \xrightarrow{f_*} & X'
 \end{array}
 =
 \begin{array}{ccc}
 X & \xrightarrow{f_*} & X' \\
 B \downarrow \cong & \cong & \downarrow B' \\
 X & \xrightarrow{f_*} & X' \\
 A \downarrow \cong & \cong & \downarrow A' \\
 X & \xrightarrow{f_*} & X'
 \end{array}$$

This comes, as the reader will easily see, to taking for  $\gamma_{z,x} : (AB)(z, x) \rightarrow (A'B')(fz, fx)$  the composite

$$\sum_{y \in X} A(z, y)B(y, x) \xrightarrow{\sum \alpha_{z,y}\beta_{y,x}} \sum_{y \in X} A'(fz, fy)B'(fy, fx) \xrightarrow{\kappa} \sum_{y' \in X'} A'(fz, y')B'(y', fx) ,$$

where the  $y$ -component of  $\kappa$  is the  $fy$ -injection into the final sum; we included the less elementary description of  $\gamma$  given above since it makes clearer the functoriality of  $M$ . The result we need is:

**3.2. LEMMA.** *The functor  $M : \mathcal{V}\text{-Gph}^{(2)} \rightarrow \mathcal{V}\text{-Gph}$  preserves filtered colimits.*

**PROOF.** Consider a functor  $(X, A, B) : \mathcal{J} \rightarrow \mathcal{V}\text{-Gph}^{(2)}$  with  $\mathcal{J}$  filtered. Using the notation above, we recall that the colimit of  $(X, A) : \mathcal{J} \rightarrow \mathcal{V}\text{-Gph}$  is  $(\bar{X}, \bar{A})$  with colimit cone  $(q_j, \alpha_j) : (X_j, A_j) \rightarrow (\bar{X}, \bar{A})$ , where  $q_j : X_j \rightarrow \bar{X}$  is the colimit cone for  $X : \mathcal{J} \rightarrow \mathbf{Set}$  and  $\hat{\alpha}_j : (q_j)_*A_j(q_j)^* \rightarrow \bar{A}$  is the colimit cone for the functor  $\tilde{A} : \mathcal{J} \rightarrow \mathcal{V}\text{-Mat}(\bar{X}, \bar{X})$  sending  $j$  to  $\tilde{A}j = (q_j)_*A_j(q_j)^*$  and sending  $\theta : j \rightarrow k$  to  $\tilde{A}\theta = (q_k)_*\hat{A}_\theta(q_k)^*$ . Similarly the colimit of  $(X, B)$  is  $(\bar{X}, \bar{B})$  with colimit cone  $(q_j, \beta_j)$ , where  $\hat{\beta}_j : (q_j)_*B_j(q_j)^* \rightarrow \bar{B}$  is the colimit cone for  $\tilde{B} : \mathcal{J} \rightarrow \mathcal{V}\text{-Mat}(\bar{X}, \bar{X})$ .



The composite of  $M$  with the functor  $(X, A, B)$  is a functor  $(X, C) : \mathcal{J} \rightarrow \mathcal{V}\text{-Gph}$  where  $C_j = A_j B_j$  and where  $C_\theta$  for  $\theta : j \rightarrow k$  is such that  $\bar{C}_\theta$  is a pasting composite of  $\bar{A}_\theta$  and  $\bar{B}_\theta$ : see the definition of  $M$  on morphisms above. This functor, of course, has the colimit cone  $(q_j, \gamma_j) : (X_j, C_j) \rightarrow (\bar{X}, \bar{C})$  where  $\hat{\gamma}_j : (q_j)_* A_j B_j (q_j)^* = (q_j)_* C_j (q_j)^* \rightarrow \bar{C}$  is the colimit cone for the functor  $\tilde{C} : \mathcal{J} \rightarrow \mathcal{V}\text{-Mat}(\bar{X}, \bar{X})$  sending  $j$  to  $\tilde{C}_j = (q_j)_* A_j B_j (q_j)^*$ .

The functor  $M$ , however, sends the colimit  $(\bar{X}, \bar{A}, \bar{B})$  of  $(X, A, B)$  to  $(\bar{X}, \bar{A}\bar{B})$ , and sends the colimit cone  $(q_j, \alpha_j, \beta_j)$  of  $(X, A, B)$  to the cone  $(q_j, \delta_j) : (X_j, A_j, B_j) \rightarrow (\bar{X}, \bar{A}\bar{B})$  where  $\delta_j$  is determined through the pasting equation

$$\begin{array}{ccc}
 X_j & \xrightarrow{(q_j)^*} & \bar{X} \\
 \downarrow A_j B_j & \xRightarrow{\bar{\delta}_j} & \downarrow \bar{A}\bar{B} \\
 X_j & \xrightarrow{(q_j)^*} & \bar{X}
 \end{array}
 =
 \begin{array}{ccc}
 X_j & \xrightarrow{(q_j)^*} & \bar{X} \\
 B_j \downarrow \xRightarrow{\bar{\beta}_j} & & \downarrow \bar{B} \\
 X_j & \xrightarrow{(q_j)^*} & \bar{X} \\
 A_j \downarrow \xRightarrow{\bar{\alpha}_j} & & \downarrow \bar{A} \\
 X_j & \xrightarrow{(q_j)^*} & \bar{X} .
 \end{array}$$

To say that  $M$  preserves the colimit of  $(X, A, B)$  is to say that the cone  $(q_j, \delta_j)$  is a colimit cone, and hence, by the above, to say that the cone

$$\hat{\delta}_j : (q_j)_* A_j B_j (q_j)^* \longrightarrow \bar{A}\bar{B}$$

is a colimit cone in  $\mathcal{V}\text{-Mat}(\bar{X}, \bar{X})$  over the functor  $\tilde{C}$ .

On the other hand, since composition of matrices commutes with colimits, the colimit cones  $\hat{\alpha}_j : \tilde{A}_j \rightarrow \bar{A}$  and  $\hat{\beta}_j : \tilde{B}_j \rightarrow \bar{B}$  give by composition a colimit cone  $\hat{\alpha}_j \hat{\beta}_k : \tilde{A}_j \tilde{B}_k \rightarrow \bar{A}\bar{B}$  over the functor  $\mathcal{J} \times \mathcal{J} \rightarrow \mathcal{V}\text{-Mat}(\bar{X}, \bar{X})$  sending  $(j, k)$  to  $\tilde{A}_j \tilde{B}_k$  and similarly defined on morphisms. Because  $\mathcal{J}$  is filtered, however, the diagonal  $\mathcal{J} \rightarrow \mathcal{J} \times \mathcal{J}$  is final; so that  $\hat{\alpha}_j \hat{\beta}_j : \tilde{A}_j \tilde{B}_j \rightarrow \bar{A}\bar{B}$  is a colimit cone for the functor  $\tilde{A}\tilde{B} : \mathcal{J} \rightarrow \mathcal{V}\text{-Mat}(\bar{X}, \bar{X})$  sending  $j$  to  $\tilde{A}_j \tilde{B}_j = (q_j)_* A_j (q_j)^* (q_j)_* B_j (q_j)^*$ .

We have the unit  $\eta_j : 1_{X_j} \rightarrow (q_j)^* (q_j)_*$  of the adjunction  $(q_j)_* \dashv (q_j)^*$ , and thus for each  $j$  a 2-cell

$$(q_j)_* A_j \eta_j B_j (q_j)^* : (q_j)_* A_j B_j (q_j)^* \rightarrow (q_j)_* A_j (q_j)^* (q_j)_* B_j (q_j)^* ,$$

which we may write as  $\zeta_j : \tilde{C}_j \rightarrow \tilde{A}_j \tilde{B}_j$ ; a straightforward calculation verifies that these are the components of a natural transformation  $\zeta : \tilde{C} \rightarrow \tilde{A}\tilde{B} : \mathcal{J} \rightarrow \mathcal{V}\text{-Mat}(\bar{X}, \bar{X})$ . Using the adjunction  $(q_j)_* \dashv (q_j)^*$  to express the  $\hat{\delta}_j$  in terms of their mates  $\bar{\delta}_j$  and hence in terms of  $\alpha$  and  $\beta$ , we find that the cone  $\hat{\delta}_j : \tilde{C}_j \rightarrow \bar{A}\bar{B}$  is just the composite of  $\zeta_j$  with the colimit cone  $\hat{\alpha}_j \hat{\beta}_j : \tilde{A}_j \tilde{B}_j \rightarrow \bar{A}\bar{B}$ . So the  $\hat{\delta}_j$  constitute a colimit cone if and only if the  $\bar{\zeta} : \bar{C} \rightarrow \bar{A}\bar{B}$  induced by  $\zeta : \tilde{C} \rightarrow \tilde{A}\tilde{B}$  is invertible.

Recall our earlier calculation of a matrix composite  $f_*Af^*$ . This gives us, for  $x, y \in \bar{X}$ ,

$$\begin{aligned} \tilde{C}_j(x, y) &= ((q_j)_*A_jB_j(q_j)^*)(x, y) \\ &= \sum_{\substack{\rho, \sigma, \tau \in X_j \\ q_j\rho=x \\ q_j\sigma=y}} A_j(\rho, \tau)B_j(\tau, \sigma) \end{aligned}$$

and

$$\begin{aligned} (\tilde{A}_j\tilde{B}_j)(x, y) &= ((q_j)_*A_j(q_j)^*(q_j)_*B_j(q_j)^*)(x, y) \\ &= \sum_{z \in \bar{X}} \sum_{\substack{r, t \in X_j \\ q_jr=x \\ q_jt=z}} \sum_{\substack{p, s \in X_j \\ q_jp=z \\ q_js=y}} A_j(r, t)B_j(p, s) ; \end{aligned}$$

and it follows easily from the explicit description of the unit  $1_{X_j} \rightarrow (q_j)^*(q_j)_*$  that  $(\zeta_j)_{x,y} : \tilde{C}_j(x, y) \rightarrow (\tilde{A}_j\tilde{B}_j)(x, y)$  is the map whose  $(\rho, \sigma, \tau)$ -component is the  $(z, r, t, p, s)$ -coprojection where  $r = \rho, s = \sigma, t = p = \tau$ , and  $z = q_j\tau$ .

We complete the proof by constructing an inverse  $\bar{\xi} : \bar{A}\bar{B} \rightarrow \bar{C}$  of  $\bar{\zeta}$ , or equally inverses  $\bar{\xi}_{x,y} : (\bar{A}\bar{B})(x, y) \rightarrow \bar{C}(x, y)$  of  $\bar{\zeta}_{x,y}$ ; here  $\bar{\xi}_{x,y}$  is to be the map induced on the colimit by a cone  $(\xi_j)_{x,y} : (\tilde{A}_j\tilde{B}_j)(x, y) \rightarrow \bar{C}(x, y)$ . By the formula above for  $(\tilde{A}_j\tilde{B}_j)(x, y)$ , it suffices to give for each  $(z, r, t, p, s)$  the appropriate component  $(\xi_j)_{x,y,z;r,t,p,s} : A_j(r, t)B_j(p, s) \rightarrow \bar{C}(x, y)$ . Now since  $q_jt = q_jp$ , there is by the filteredness of  $\mathcal{J}$  some  $\theta : j \rightarrow k$  with  $X_\theta t = X_\theta p = t' \in X_k$ , say. Write  $r'$  for  $X_\theta r$  and  $s'$  for  $X_\theta s$ . We take for  $(\xi_j)_{x,y,z;r,t,p,s}$  the composite

$$A_j(r, t)B_j(p, s) \xrightarrow{(A_\theta)_{r,t}(B_\theta)_{p,s}} A_k(r', t')B_k(t', s') \xrightarrow{\lambda} \tilde{C}_k(x, y) \xrightarrow{(\hat{\gamma}_k)_{x,y}} \bar{C}(x, y) ,$$

where  $\lambda$  is the appropriate coprojection in the expression above for  $\tilde{C}_j(x, y)$ , but now with  $k$  in place of  $j$ . It is easy to verify, first, that  $(\xi_j)_{x,y,z;r,t,p,s}$  is independent of our choice of a  $\theta : j \rightarrow k$  with  $X_\theta t = X_\theta p$ , so that  $(\xi_j)_{x,y}$  is well-defined; and second that the  $(\xi_j)_{x,y} : (\tilde{A}_j\tilde{B}_j)(x, y) \rightarrow \bar{C}(x, y)$  constitute a cone, thus inducing a map  $\bar{\xi}_{x,y} : \bar{A}\bar{B}(x, y) \rightarrow \bar{C}(x, y)$  determined by  $\bar{\xi}_{x,y}(\hat{\alpha}_j\hat{\beta}_j)_{x,y} = (\xi_j)_{x,y}$ .

That  $\bar{\xi}_{x,y}\bar{\zeta}_{x,y} = 1$  follows easily because, in applying  $\bar{\xi}_{x,y}$  on the image of  $\bar{\zeta}_{x,y}$  we may, since here  $t = p = \tau$ , take  $\theta : j \rightarrow k$  to be  $1_j$ . To say that  $\bar{\zeta}_{x,y}\bar{\xi}_{x,y} = 1$  is to say that  $\bar{\zeta}_{x,y}\bar{\xi}_{x,y}(\hat{\alpha}_j\hat{\beta}_j)_{x,y} = (\hat{\alpha}_j\hat{\beta}_j)_{x,y}$  for each  $j$ . However  $\bar{\zeta}_{x,y}\bar{\xi}_{x,y}(\hat{\alpha}_j\hat{\beta}_j)_{x,y} = \bar{\zeta}_{x,y}(\xi_j)_{x,y}$ , whose  $(z; r, t, p, s)$ -component by the above is

$$\bar{\zeta}_{x,y}(\hat{\gamma}_k)_{x,y}\lambda((A_\theta)_{r,t}(B_\theta)_{p,s}) = (\hat{\alpha}_k\hat{\beta}_k)_{x,y}(\zeta_k)_{x,y}\lambda((A_\theta)_{r,t}(B_\theta)_{p,s}) ,$$

and it follows from the explicit description above of  $(\zeta_k)_{x,y}$  that  $(\zeta_k)_{x,y}\lambda$  is just the coprojection  $A_k(r', t')B_k(t', s') \rightarrow (\tilde{A}_k\tilde{B}_k)(x, y)$ , which we shall write as  $\kappa_k$ . If we similarly

write  $\kappa_j$  for the coprojection  $A_j(r, t)B_j(p, s) \rightarrow (\tilde{A}_j \tilde{B}_j)(x, y)$ , we have  $\kappa_k((A_\theta)_{r,t}(B_\theta)_{p,s}) = (\tilde{A}_\theta \tilde{B}_\theta)_{x,y} \kappa_j$ , so that  $(\hat{\alpha}_k \hat{\beta}_k)_{x,y} \kappa_k((A_\theta)_{r,t}(B_\theta)_{p,s}) = (\hat{\alpha}_k \hat{\beta}_k)_{x,y} (\tilde{A}_\theta \tilde{B}_\theta)_{x,y} \kappa_j = (\hat{\alpha}_j \hat{\beta}_j)_{x,y} \kappa_j$ , which is the  $(z; r, t, p, s)$ -component of  $(\hat{\alpha}_j \hat{\beta}_j)_{x,y}$ , as desired. So the  $\tilde{\zeta}_{x,y}$  are indeed invertible, which completes the proof.  $\blacksquare$

We shall now describe the endofunctor  $T$  of  $\mathcal{V}$ -**Gph** underlying the “free  $\mathcal{V}$ -category” monad. Recall from [2] that  $T$  sends a  $\mathcal{V}$ -graph  $(X, A)$  to  $(X, A')$  where  $A' = \sum_{n \in \mathbb{N}} A^n$  is the free monoid on  $A$  in the monoidal category given by  $\mathcal{V}$ -**Mat** $(X, X)$  with matrix multiplication as its tensor product; and that the unit  $(X, A) \rightarrow (X, A')$  of the adjunction is  $(1, \rho_A)$  where  $\rho_A : A \rightarrow A'$  is the injection of the summand  $A = A^1$  into  $\sum A^n$ . From this we can calculate the value of  $T$  on morphisms, which leads to the following description of  $T$ . For each  $n \in \mathbb{N}$  there is an endofunctor  $T_n$  of  $\mathcal{V}$ -**Gph** sending  $(X, A)$  to  $(X, A^n)$ ; and because  $PT_n = P$ , these  $T_n$  are the components of a functor  $T_{\mathbb{N}} : \mathcal{V}\text{-Gph} \rightarrow \mathcal{V}\text{-Gph}^{(\mathbb{N})}$ ; whereupon  $T$  is the composite  $ST_{\mathbb{N}}$ , where  $S : \mathcal{V}\text{-Gph}^{(\mathbb{N})} \rightarrow \mathcal{V}\text{-Gph}$  is the functor so denoted in Lemma 3.1. Since  $S$  preserves all colimits by Lemma 3.1,  $T$  will be finitary (that is, will preserve filtered colimits) if  $T_{\mathbb{N}}$  is so. Since the projections  $\mathcal{V}\text{-Gph}^{(\mathbb{N})} \rightarrow \mathcal{V}\text{-Gph}$  jointly create colimits,  $T_{\mathbb{N}}$  will be finitary if each  $T_n$  is so. However  $T_1$  is the identity endofunctor 1 of  $\mathcal{V}\text{-Gph}$ , while  $T_2$  is the composite  $M(1, 1)$ , where  $(1, 1) : \mathcal{V}\text{-Gph} \rightarrow \mathcal{V}\text{-Gph}^{(2)}$  is the functor each of whose components is 1; and  $T_{n+1}$  for  $n \geq 1$  is (isomorphic to) the composite  $M(T_n, 1)$ . Since the projections  $\mathcal{V}\text{-Gph}^{(2)} \rightarrow \mathcal{V}\text{-Gph}$  jointly create colimits, it follows inductively from Lemma 3.2 that  $T_n$  is finitary for  $n \geq 1$ .

It remains to consider the endofunctor  $T_0$  of  $\mathcal{V}\text{-Gph}$  sending  $(X, A)$  to  $(X, 1_X)$ , where  $1_X$  is the identity matrix with  $(1_X)_{x,y}$  being  $I$  for  $x = y$  and 0 otherwise. This is the composite of the forgetful functor  $P : \mathcal{V}\text{-Gph} \rightarrow \mathbf{Set}$  and the evident functor  $H : \mathbf{Set} \rightarrow \mathcal{V}\text{-Gph}$  sending  $X$  to  $(X, 1_X)$ . Since  $P$  preserves all colimits, it will suffice to show that  $H$  preserves filtered colimits. Suppose then that  $X : \mathcal{J} \rightarrow \mathbf{Set}$  with  $\mathcal{J}$  filtered has as before the colimit cone  $(q_j : X_j \rightarrow \bar{X})$ , and consider the colimit of  $HX$ ; our claim is that the colimit of the  $(X_j, 1_{X_j})$  is  $(\bar{X}, 1_{\bar{X}})$ . By our description of colimits in  $\mathcal{V}\text{-Gph}$ , we have to show that  $1_{\bar{X}}$  is the colimit in  $\mathcal{V}\text{-Mat}(\bar{X}, \bar{X})$  of the  $(q_j)_* 1_{X_j} (q_j)^*$ . Since

$$((q_j)_* 1_{X_j} (q_j)^*)(x, y) = \sum_{\substack{q_j r = x \\ q_j s = y}} (1_{X_j})(r, s) ,$$

there is nothing to prove for  $x \neq y$ , the cone being constant at 0. For  $x = y$  the above gives

$$((q_j)_* 1_{X_j} (q_j)^*)(x, x) = q_j^{-1}(x) \bullet I ,$$

the coproduct of  $q_j^{-1}(x)$  copies of  $I$ ; and we are claiming that the colimit in  $\mathcal{V}$  of the  $q_j^{-1}(x) \bullet I$  is  $I$ . However  $(\ ) \bullet I : \mathbf{Set} \rightarrow \mathcal{V}$  preserves colimits, so that it suffices to observe that in  $\mathbf{Set}$  we have  $\text{colim}(q_j^{-1}(x)) = 1$ . But filtered colimits in  $\mathbf{Set}$  commute with finite limits; and the above is precisely what we get on pulling back the colimit  $q_j : X_j \rightarrow \bar{X}$  along  $x : 1 \rightarrow \bar{X}$ . This completes the proof of:

3.3. THEOREM. *The monad on  $\mathcal{V}\text{-Gph}$  whose algebras are  $\mathcal{V}$ -categories is finitary.*

An equivalent formulation is:

3.4. COROLLARY. *The forgetful functor  $U : \mathcal{V}\text{-Cat} \rightarrow \mathcal{V}\text{-Gph}$  is finitary.*

4.  $\mathcal{V}\text{-Cat}$  is locally presentable if  $\mathcal{V}$  is so

As in Section 3, we continue to suppose that the monoidal category  $\mathcal{V}$  is cocomplete and that the functors  $A \otimes -$  and  $- \otimes A$  preserve colimits, as they surely do when  $\mathcal{V}$  is closed. To avoid pathologies in our use of the “strong generator” notion, we further suppose that  $\mathcal{V}_0$  admits arbitrary cointersections of strong epimorphisms, which ensures that (strong epimorphisms, monomorphisms) is a factorization system on  $\mathcal{V}_0$ . This presents no problem, since our main goal is the study of the case where  $\mathcal{V}_0$  is locally presentable.

It is convenient to introduce, for each object  $G$  of  $\mathcal{V}$ , the  $\mathcal{V}$ -graph  $(2, \bar{G})$  having  $2 = \{0, 1\}$  for its set of objects and having

$$\bar{G}(0, 1) = G, \quad \bar{G}(0, 0) = \bar{G}(1, 1) = \bar{G}(1, 0) = 0,$$

where this last 0 is the initial object of  $\mathcal{V}$ ; that is to say,  $\bar{G}$  is the 2-by-2 matrix  $\begin{pmatrix} 0 & G \\ 0 & 0 \end{pmatrix}$ .

To give a morphism  $(2, \bar{G}) \rightarrow (X, A)$  of  $\mathcal{V}$ -graphs is just to give a pair  $x, y \in X$  and a morphism  $u : G \rightarrow A(x, y)$  in  $\mathcal{V}$ .

The forgetful functor  $P : \mathcal{V}\text{-Gph} \rightarrow \mathbf{Set}$  sending the  $\mathcal{V}$ -graph  $(X, A)$  to  $X$  clearly has a left adjoint  $D$  sending the set  $X$  to the  $\mathcal{V}$ -graph  $(X, 0_X)$ , where  $0_X$  is the initial object of  $\mathcal{V}\text{-Mat}(X, X)$  given by  $0_X(x, x') = 0$ .

4.1. LEMMA. *A morphism  $(f, \alpha) : (X, A) \rightarrow (Y, B)$  in  $\mathcal{V}\text{-Gph}$  is monomorphic if and only if  $f : X \rightarrow Y$  is an injective function and each  $\alpha_{x,x'} : A(x, x') \rightarrow B(fx, fx')$  is a monomorphism in  $\mathcal{V}$  (that is, in  $\mathcal{V}_0$ ).*

PROOF. The “if” part being clear from the definition of composition in  $\mathcal{V}\text{-Gph}$ , it suffices to prove the “only if” part; so suppose that  $(f, \alpha)$  is monomorphic in  $\mathcal{V}\text{-Gph}$ . Then  $f$  is injective because  $P : \mathcal{V}\text{-Gph} \rightarrow \mathbf{Set}$ , having a left adjoint, preserves monomorphisms. Suppose that, for some  $x, x' \in X$ , maps  $\beta, \gamma : G \rightarrow A(x, x')$  in  $\mathcal{V}$  satisfy  $\alpha_{x,x'}\beta = \alpha_{x,x'}\gamma$ , and define  $g : 2 \rightarrow X$  by setting  $g0 = x$  and  $g1 = x'$ ; now the morphisms  $(g, \beta), (g, \gamma) : (2, \bar{G}) \rightarrow (X, A)$  have the same composite with  $(f, \alpha) : (X, A) \rightarrow (Y, B)$ , whence  $\beta = \gamma$ . Thus  $\alpha_{x,x'}$  is indeed monomorphic. ■

4.2. LEMMA. *If a set  $\mathcal{G}$  of objects constitutes a strong generator of  $\mathcal{V}_0$ , then the set  $\{(2, \bar{G}) \mid G \in \mathcal{G} \text{ or } G = 0\}$  constitutes a strong generator of  $\mathcal{V}\text{-Gph}$ .*

PROOF. We prove the assertion in the equivalent form — see Section 2 above — that the totality of maps in  $\mathcal{V}\text{-Gph}$  into the object  $(Y, B)$  having domain one of the  $(2, \bar{G})$  with  $G \in \mathcal{G} \cup \{0\}$  factorizes through no proper subobject of  $(Y, B)$  and is therefore jointly a strong epimorphism. Suppose then that  $(f, \alpha) : (X, A) \rightarrow (Y, B)$  is a monomorphism in

$\mathcal{V}$ -**Gph** through which every  $(g, \beta) : (2, \bar{G}) \rightarrow (Y, B)$  with  $G \in \mathcal{G} \cup \{0\}$  factorizes. To give a map from  $(2, \bar{0}) = (2, 0_2)$  into  $(Y, B)$  is just to give two elements of  $Y$ ; and since every such map factorizes through  $(f, \alpha)$ , the injection  $f$  is in fact a bijection. Since every  $(g, \beta) : (2, \bar{G}) \rightarrow (Y, B)$  with  $G \in \mathcal{G}$  factorizes through  $(f, \alpha)$ , every map  $G \rightarrow B(fx, fx')$  in  $\mathcal{V}$  factorizes through the monomorphism  $\alpha_{x,x'} : A(x, x') \rightarrow B(fx, fx')$ , which is therefore invertible, because  $\mathcal{G}$  is a strong generator for  $\mathcal{V}_0$ . Thus the monomorphism  $(f, \alpha)$  is indeed invertible.  $\blacksquare$

We now examine the “presentability” of such a strong generator for  $\mathcal{V}$ -**Gph**. Note that  $\mathcal{V}_0(0, -) : \mathcal{V}_0 \rightarrow \mathbf{Set}$  is the functor constant at 1, which preserves all connected colimits; so that the object 0 of  $\mathcal{V}_0$  is  $\lambda$ -presentable for any regular cardinal  $\lambda$ .

4.3. LEMMA. *If, for some regular cardinal  $\lambda$ , the object  $G$  is  $\lambda$ -presentable in  $\mathcal{V}_0$ , then  $(2, \bar{G})$  is  $\lambda$ -presentable in  $\mathcal{V}$ -**Gph**.*

PROOF. Consider as in Section 3 above the colimit cone  $(q_j, \alpha_j) : (X_j, A_j) \rightarrow (\bar{X}, \bar{A})$  of a functor  $(X, A) : \mathcal{J} \rightarrow \mathcal{V}$ -**Gph**, where the category  $\mathcal{J}$  is  $\lambda$ -filtered; we are to show that the functor  $\mathcal{V}$ -**Gph** $((2, \bar{G}), -) : \mathcal{V}$ -**Gph**  $\rightarrow$  **Set** preserves every such colimit; equivalently, we are to prove bijective the canonical comparison

$$\kappa : \text{colim}_{k \in \mathcal{J}} \mathcal{V}\text{-Gph}((2, \bar{G}), (X_k, A_k)) \rightarrow \mathcal{V}\text{-Gph}((2, \bar{G}), (\bar{X}, \bar{A}))$$

of sets. We begin by proving  $\kappa$  surjective; that is to say, that every map  $(g, \tau) : (2, \bar{G}) \rightarrow (\bar{X}, \bar{A})$  factorizes through some  $(q_k, \alpha_k) : (X_k, A_k) \rightarrow (\bar{X}, \bar{A})$ . To give  $(g, \tau)$  is to give a function  $g : 2 \rightarrow \bar{X}$  picking out elements  $x, y \in \bar{X}$  and to give a map  $\tau : G \rightarrow \bar{A}(x, y)$  in  $\mathcal{V}$ . We recall from Section 3, however, that the  $(\hat{\alpha}_j)_{x,y} : \tilde{A}_j(x, y) \rightarrow \bar{A}(x, y)$  constitute a colimit cone for the functor  $\tilde{A}(x, y) : \mathcal{J} \rightarrow \mathcal{V}$ ; so,  $G$  being  $\lambda$ -presentable in  $\mathcal{V}_0$ , the map  $\tau : G \rightarrow \bar{A}(x, y)$  factorizes as

$$G \xrightarrow{\sigma} \tilde{A}_j(x, y) \xrightarrow{(\hat{\alpha}_j)_{x,y}} \bar{A}(x, y)$$

for some  $j \in \mathcal{J}$ . Here  $\tilde{A}_j$  is the object  $(q_j)_* A_j (q_j)^*$  of  $\mathcal{V}\text{-Mat}(\bar{X}, \bar{X})$ , so that

$$\tilde{A}_j(x, y) = \sum_{\substack{q_j t = x \\ q_j s = y}} A_j(t, s).$$

This coproduct, however, is the  $\lambda$ -filtered colimit of its sub-coproducts indexed by subsets of  $q_j^{-1}(x) \times q_j^{-1}(y)$  of cardinality less than  $\lambda$ ; so,  $G$  being  $\lambda$ -presentable,  $\sigma$  factorizes through such a sub-coproduct, say  $\sum_{\nu \in N} A_j(t_\nu, s_\nu)$  where  $q_j t_\nu = x$  and  $q_j s_\nu = y$  for all  $\nu \in N$  and where  $\text{card} N < \lambda$ . Using yet again the  $\lambda$ -filteredness of  $\mathcal{J}$ , there is some arrow  $\theta : j \rightarrow k$  in  $\mathcal{J}$  for which all the  $X_\theta t_\nu$  are equal and all the  $X_\theta s_\nu$  are equal: say

$$X_\theta t_\nu = \bar{t} \in X_k \text{ and } X_\theta s_\nu = \bar{s} \in X_k \text{ for all } \nu \in N.$$

Using the factorization above of  $\tau$ , we have

$$\tau = (\hat{\alpha}_j)_{x,y} \sigma = (\hat{\alpha}_k)_{x,y} (\tilde{A}_\theta)_{x,y} \sigma = (\hat{\alpha}_k)_{x,y} \rho,$$

where  $\rho = (\tilde{A}_\theta)_{x,y}\sigma$ , which by Section 3 is in fact  $\left( (q_k)_* \widehat{A}_\theta (q_k)^* \right)_{x,y} \sigma$ . The point is that this map

$$\rho : G \rightarrow \tilde{A}_k(x, y) = \sum_{\substack{q_k t' = x \\ q_k s' = y}} A_k(t', s')$$

factorizes through the coprojection of a single summand  $A_k(\bar{t}, \bar{s})$ , say as the composite of this coprojection with the map  $\varphi : G \rightarrow A_k(\bar{t}, \bar{s})$  of  $\mathcal{V}$ . Now the pair  $(\bar{t}, \bar{s})$  determines a function  $h : 2 \rightarrow X_k$ , so that we have in  $\mathcal{V}\text{-Gph}$  the map  $(h, \varphi) : (2, \bar{G}) \rightarrow (X_k, A_k)$ ; but the composite of this with  $(q_k, \alpha_k)$  is  $(g, \tau)$ . So  $(g, \tau)$  does indeed factorize through some  $(q_k, \alpha_k)$ , which completes the proof that the canonical comparison  $\kappa$  is surjective.

It remains to show that  $\kappa$  is injective. Suppose then that  $(h, \varphi) : (2, \bar{G}) \rightarrow (X_k, A_k)$  and  $(h', \varphi') : (2, \bar{G}) \rightarrow (X_{k'}, A_{k'})$  are maps in  $\mathcal{V}\text{-Gph}$  with  $(q_k, \alpha_k)(h, \varphi) = (q_{k'}, \alpha_{k'})(h', \varphi')$ . We are to show that there exist  $\theta : k \rightarrow j$  and  $\theta' : k' \rightarrow j$  in  $\mathcal{J}$ , with  $(X_\theta, \tilde{A}_\theta)(h, \varphi) = (X_{\theta'}, \tilde{A}_{\theta'})(h', \varphi')$ . Since  $\mathcal{J}$  is  $\lambda$ -filtered, there certainly do exist maps  $\theta : k \rightarrow j$  and  $\theta'_0 : k' \rightarrow j$ , and so without loss of generality we may suppose that  $k = k'$ .

Write  $(t, s)$  for  $(h0, h1)$  and  $(t', s')$  for  $(h'0, h'1)$ , so that  $\varphi : G \rightarrow A_k(t, s)$  and  $\varphi' : G \rightarrow A_k(t', s')$ . Since  $q_k h = q_k h'$ , there is some  $\theta : k \rightarrow j$  with  $X_\theta t = X_\theta t'$  and  $X_\theta s = X_\theta s'$ ; in other words, we may suppose without loss of generality that  $t = t'$  and  $s = s'$ . Now, therefore,  $h = h' : 2 \rightarrow X_k$  corresponds to  $(t, s) \in X_k$ , and  $\varphi, \varphi' : G \rightarrow A_k(t, s)$  have  $\alpha_k \varphi = \alpha_k \varphi'$ . Write  $\bar{t}$  for  $q_k t \in \bar{X}$ , and  $\bar{s}$  for  $q_k s$ , and recall from Section 3 that we have in  $\mathcal{V}$  the colimit cone  $((\hat{\alpha}_j)_{\bar{t}, \bar{s}} : \tilde{A}_j(\bar{t}, \bar{s}) \rightarrow \bar{A}(\bar{t}, \bar{s}))_{j \in \mathcal{J}}$ , where

$$\tilde{A}_j(\bar{t}, \bar{s}) = \sum_{\substack{q_j \tilde{t} = \bar{t} \\ q_j \tilde{s} = \bar{s}}} A_j(\tilde{t}, \tilde{s}).$$

Composing  $\varphi$  and  $\varphi'$  with the  $(t, s)$ -coprojection for  $\tilde{A}_k$  gives us two maps  $\psi, \psi' : G \rightarrow \tilde{A}_k(\bar{t}, \bar{s})$  in  $\mathcal{V}$  with  $(\hat{\alpha}_k)_{\bar{t}, \bar{s}} \psi = (\hat{\alpha}_k)_{\bar{t}, \bar{s}} \psi'$ . Because  $\mathcal{J}$  is  $\lambda$ -filtered and  $G$  is  $\lambda$ -presentable in  $\mathcal{V}_0$ , there is some  $\theta : k \rightarrow j$  in  $\mathcal{J}$  for which  $\tilde{A}_\theta(\bar{t}, \bar{s})\psi = \tilde{A}_\theta(\bar{t}, \bar{s})\psi'$ . When we recall the definition of the functor  $\tilde{A}$ , we see that this gives exactly the equality  $(X_\theta, A_\theta)(h, \varphi) = (X_\theta, A_\theta)(h, \varphi')$  that we need for the injectivity of  $\kappa$ . ■

Since, as we have remarked, the object 0 of  $\mathcal{V}_0$  is  $\lambda$ -presentable for any regular cardinal  $\lambda$ , it follows from Lemmas 4.2 and 4.3 that, under the standing hypotheses of this section:

4.4. PROPOSITION.  $\mathcal{V}\text{-Gph}$  is locally  $\lambda$ -presentable when  $\mathcal{V}_0$  is so. ■

Combining this with Theorem 3.3 and using [5, Satz 10.3], we conclude that :

4.5. THEOREM. If  $\mathcal{V}$  is a monoidal closed category whose underlying ordinary category  $\mathcal{V}_0$  is locally  $\lambda$ -presentable, then  $\mathcal{V}\text{-Cat}$  is also locally  $\lambda$ -presentable.

We can be more specific, in the sense of actually exhibiting a strong generator for  $\mathcal{V}\text{-Cat}$  consisting of  $\lambda$ -presentable objects. We make use of the following simple and well-known general observations:

4.6. LEMMA. *Let  $F \dashv U : \mathcal{A} \rightarrow \mathcal{B}$  where  $\mathcal{A}$  and  $\mathcal{B}$  are cocomplete categories. Then (i) for an object  $G$  of  $\mathcal{B}$ , the object  $FG$  of  $\mathcal{A}$  is  $\lambda$ -presentable if  $G$  is  $\lambda$ -presentable and  $U$  preserves  $\lambda$ -filtered colimits; and (ii) if a small set  $\mathcal{G}$  of objects of  $\mathcal{B}$  constitutes a strong generator of  $\mathcal{B}$ , and if  $U$  reflects isomorphisms (as it surely does whenever it is monadic), then the set  $\{FG \mid G \in \mathcal{G}\}$  constitutes a strong generator of  $\mathcal{A}$ .*

Let us use  $F$  now for the left adjoint of the forgetful  $U : \mathcal{V}\text{-Cat} \rightarrow \mathcal{V}\text{-Gph}$ . When  $\mathcal{V}_0$  is locally  $\lambda$ -presentable, we can take for  $\mathcal{G}$  the full subcategory  $\mathcal{V}_\lambda$  of  $\mathcal{V}_0$  given by the  $\lambda$ -presentable objects, noting that it contains the initial object  $0$ . By Lemmas 4.2 and 4.3, the  $(2, \bar{G})$  for  $G \in \mathcal{V}_\lambda$  constitute a strong generator of  $\mathcal{V}\text{-Gph}$  consisting of  $\lambda$ -presentable objects. By Lemma 4.6 and Corollary 3.4, therefore, the  $F(2, \bar{G})$  for  $G \in \mathcal{V}_\lambda$  constitute a strong generator of  $\mathcal{V}\text{-Cat}$  consisting of  $\lambda$ -presentable objects. In future we shall write  $\mathfrak{2}_G$  for the  $\mathcal{V}$ -category  $F(2, \bar{G})$ ; it is characterized by the observation that to give a  $\mathcal{V}$ -functor from  $\mathfrak{2}_G$  to a  $\mathcal{V}$ -category  $B$  is to give objects  $x$  and  $y$  of  $B$  along with a map  $G \rightarrow B(x, y)$  in  $\mathcal{V}$ . The reader will easily verify that  $\mathfrak{2}_G$  has two objects  $0$  and  $1$ , with  $\mathfrak{2}_G(0, 0) = \mathfrak{2}_G(1, 1) = I$ ,  $\mathfrak{2}_G(0, 1) = G$ ,  $\mathfrak{2}_G(1, 0) = 0$ , and with the evident composition.

A standard result from the theory of locally presentable categories now gives:

4.7. PROPOSITION. *When  $\mathcal{V}_0$  is locally  $\lambda$ -presentable, the class of  $\lambda$ -presentable objects in  $\mathcal{V}\text{-Cat}$  is the closure in  $\mathcal{V}\text{-Cat}$  under  $\lambda$ -small colimits of the  $\mathcal{V}$ -categories  $\mathfrak{2}_G$ , where  $G$  is a  $\lambda$ -presentable object of  $\mathcal{V}$ . ■*

Recall from Section 2 above that a monoidal closed category  $\mathcal{V}$  is locally  $\lambda$ -presentable as a closed category when its underlying ordinary category  $\mathcal{V}_0$  is locally  $\lambda$ -presentable, and the  $\lambda$ -presentable objects of  $\mathcal{V}_0$  are closed under the monoidal structure. Although our interest in local presentability for closed categories is rather secondary, we nonetheless record:

4.8. PROPOSITION. *If the symmetric monoidal closed category  $\mathcal{V}$  is locally  $\lambda$ -presentable as a closed category, then so is  $\mathcal{V}\text{-Cat}$ .*

PROOF. We must show that the  $\lambda$ -presentable  $\mathcal{V}$ -categories are closed under tensor product. By [9, (5.2)] it suffices to show that  $\mathfrak{2}_G \otimes \mathfrak{2}_H$  is  $\lambda$ -presentable for all  $G, H \in \mathcal{V}_\lambda$ .

Write  $\mathcal{I}$  for the  $\mathcal{V}$ -category with a single object  $*$  and  $\mathcal{I}(*, *) = I$ ; to give a  $\mathcal{V}$ -functor  $\mathcal{I} \rightarrow \mathcal{A}$  is just to give an object of  $\mathcal{A}$ . Thus  $\mathcal{V}\text{-Cat}(\mathcal{I}, -) : \mathcal{V}\text{-Cat} \rightarrow \mathbf{Set}$  is the functor sending a  $\mathcal{V}$ -category to its set of objects. This has a right adjoint, and so preserves all colimits, whence  $\mathcal{I}$  is certainly  $\lambda$ -presentable.

The  $\mathcal{V}$ -category  $\mathcal{C} = \mathfrak{2}_G \otimes \mathfrak{2}_H$  has four objects:  $(0, 0)$ ,  $(0, 1)$ ,  $(1, 0)$ , and  $(1, 1)$ , and hom-objects

$$\mathcal{C}((i, j), (i', j')) = \begin{cases} G & \text{if } i = 0, i' = 1, j = j' \\ H & \text{if } i = i', j = 0, j' = 1 \\ G \otimes H & \text{if } i = i' = 0, j = j' = 1 \\ I & \text{if } i = i', j = j' \\ 0 & \text{otherwise} \end{cases}$$

with the obvious composition maps. To give a  $\mathcal{V}$ -functor  $T : \mathfrak{2}_G \otimes \mathfrak{2}_H \rightarrow \mathcal{A}$ , therefore, is to give four objects  $A = S(0, 0)$ ,  $B = S(0, 1)$ ,  $C = S(1, 0)$ ,  $D = S(1, 1)$  of  $\mathcal{A}$ , along with maps  $\alpha : G \rightarrow \mathcal{A}(A, C)$ ,  $\beta : G \rightarrow \mathcal{A}(B, D)$ ,  $\gamma : H \rightarrow \mathcal{A}(A, B)$ , and  $\delta : H \rightarrow \mathcal{A}(C, D)$  in  $\mathcal{V}$  rendering commutative the diagram

$$\begin{array}{ccc}
 G \otimes H & \xrightarrow{\beta \otimes \gamma} & \mathcal{A}(B, D) \otimes \mathcal{A}(A, B) \\
 \tau \downarrow & & \downarrow M \\
 H \otimes G & \xrightarrow{\delta \otimes \alpha} \mathcal{A}(C, D) \otimes \mathcal{A}(A, C) \xrightarrow{M} & \mathcal{A}(A, D)
 \end{array} \tag{*}$$

wherein  $\tau$  denotes the symmetry isomorphism and  $M$  the composition maps in  $\mathcal{A}$ .

To give  $A, C, D$  along with  $\alpha$  and  $\delta$  is to give a  $\mathcal{V}$ -functor  $R : \mathfrak{3}_{G,H} \rightarrow \mathcal{A}$ , where  $\mathfrak{3}_{G,H}$  is the pushout

$$\begin{array}{ccc}
 \mathcal{I} & \xrightarrow{0} & \mathfrak{2}_H \\
 1 \downarrow & & \downarrow \\
 \mathfrak{2}_G & \longrightarrow & \mathfrak{3}_{G,H}
 \end{array}$$

in  $\mathcal{V}\text{-Cat}$ . Similarly to give  $A, B, D$  along with  $\beta$  and  $\gamma$  is to give a  $\mathcal{V}$ -functor  $S : \mathfrak{3}_{H,G} \rightarrow \mathcal{A}$ . To give  $T : \mathfrak{2}_G \otimes \mathfrak{2}_H \rightarrow \mathcal{A}$ , therefore, is to give  $R$  and  $S$  with the same  $A$  and  $D$  and satisfying (\*); which is to say that  $\mathfrak{2}_G \otimes \mathfrak{2}_H$  is the pushout

$$\begin{array}{ccc}
 \mathfrak{2}_{G \otimes H} & \xrightarrow{N} & \mathfrak{3}_{H,G} \\
 M \downarrow & & \downarrow \\
 \mathfrak{3}_{G,H} & \longrightarrow & \mathfrak{2}_G \otimes \mathfrak{2}_H
 \end{array}$$

in  $\mathcal{V}\text{-Cat}$ , where  $M$  and  $N$  are the evident  $\mathcal{V}$ -functors. Since  $\mathcal{I}$ ,  $\mathfrak{2}_G$ , and  $\mathfrak{2}_H$  are  $\lambda$ -presentable, and since the  $\lambda$ -presentables are closed under finite colimits, it follows that  $\mathfrak{2}_G \otimes \mathfrak{2}_H$  is  $\lambda$ -presentable, as desired. ■

### 5. $\mathcal{V}\text{-Cat}$ is locally bounded if $\mathcal{V}$ is so

For the first part of this section we suppose only that  $\mathcal{V}$  is cocomplete and monoidal closed. The functor  $\text{ob} : \mathcal{V}\text{-Cat} \rightarrow \mathbf{Set}$  sending a  $\mathcal{V}$ -category to its set of objects then has both adjoints: the left adjoint  $D$  sends a set  $X$  to the “discrete”  $\mathcal{V}$ -category with object-set  $X$  and  $DX(x, y)$  equal to 0 unless  $x = y$  in which case it is the unit  $I$ , and the right adjoint  $C$  sends  $X$  to the “chaotic”  $\mathcal{V}$ -category with object-set  $X$  and  $CX(x, y) = 1$ .

For a  $\mathcal{V}$ -functor  $F : \mathcal{A} \rightarrow \mathcal{B}$ , we consider the set  $K$  of pairs  $(A, B)$  of objects of  $\mathcal{A}$  with  $FA = FB$ , and form the coequalizer  $Q : \mathcal{A} \rightarrow \mathcal{C}$  in  $\mathcal{V}\text{-Cat}$  of the two “projections”  $P_1, P_2 : DK \rightarrow \mathcal{A}$ . Since  $FP_1 = FP_2$ , there is a unique  $\mathcal{V}$ -functor  $I : \mathcal{C} \rightarrow \mathcal{B}$  satisfying  $IQ = F$ . Clearly  $Q$  is invertible if and only if  $F$  is injective on objects; if  $I$  is invertible then we say that  $F$  is a *quotient on objects*. Since the “congruence”  $K$  arising from  $F$



is the same as that arising from  $Q$ , we see that  $Q$  is a quotient on objects, while  $I$  is injective on objects by construction. The factorization is clearly functorial, and so we obtain a factorization system  $(\mathcal{Q}, \mathcal{I})$  on  $\mathcal{V}\text{-Cat}$  in which  $\mathcal{Q}$  consists of the quotients on objects, and  $\mathcal{I}$  consists of those  $\mathcal{V}$ -functors that are injective on objects. (Although each  $Q \in \mathcal{Q}$  is an epimorphism in  $\mathcal{V}\text{-Cat}$ —in fact a regular one—there can be  $\mathcal{V}$ -functors which are injective on objects but not monomorphic, and so  $(\mathcal{Q}, \mathcal{I})$  is not proper.)

Before turning to the main results of the section, recall that if  $\mathcal{H}$  is a class of arrows in  $\mathcal{V}$ , a  $\mathcal{V}$ -functor  $F : \mathcal{A} \rightarrow \mathcal{B}$  is said to be *locally in  $\mathcal{H}$*  if each  $F : \mathcal{A}(A, B) \rightarrow \mathcal{B}(FA, FB)$  is in  $\mathcal{H}$ .

We now suppose that  $\mathcal{V}$  is locally  $\lambda$ -bounded as a closed category, with respect to the proper factorization system  $(\mathcal{E}, \mathcal{M})$ .

Write  $\mathcal{M}'$  for the class of  $\mathcal{V}$ -functors which are injective on objects and locally in  $\mathcal{M}$ ; clearly every such  $\mathcal{V}$ -functor is a monomorphism. Because of the  $(\mathcal{E}, \mathcal{M})$ -generator  $\mathcal{G}$ , an object of  $\mathcal{V}$  has only a small set of  $\mathcal{M}$ -subobjects, from which it follows that an object of  $\mathcal{V}\text{-Cat}$  has only a small set of  $\mathcal{M}'$ -subobjects, and therefore admits arbitrary intersections of  $\mathcal{M}'$ -subobjects. Since  $\mathcal{M}'$  is clearly closed under composition and intersections, and stable under pullback, it forms, by [3, Lemma 3.1], part of a factorization system  $(\mathcal{E}', \mathcal{M}')$  on  $\mathcal{V}\text{-Cat}$ . Since  $(\mathcal{E}, \mathcal{M})$  is proper, every coretraction in  $\mathcal{V}$  lies in  $\mathcal{M}$ , and one now easily shows that every coretraction in  $\mathcal{V}\text{-Cat}$  lies in  $\mathcal{M}'$ , and so that  $(\mathcal{E}', \mathcal{M}')$  is proper.

It takes a little work to compute  $\mathcal{E}'$ , although it is easy to see that a  $\mathcal{V}$ -functor  $F : \mathcal{A} \rightarrow \mathcal{B}$  in  $\mathcal{E}'$  must be surjective on objects, since otherwise it would factorize through some non-invertible  $J : \mathcal{C} \rightarrow \mathcal{B}$  which is injective on objects and fully faithful, and therefore lies in  $\mathcal{M}'$ . Now consider, for an  $F : \mathcal{A} \rightarrow \mathcal{B}$  that *is* surjective on objects, its  $(\mathcal{Q}, \mathcal{I})$ -factorization  $F = IQ$ . Since here  $I$ , like  $F$ , is surjective on objects, it is in fact bijective on objects. The quotient-on-objects  $\mathcal{V}$ -functor  $Q$ , being a regular epimorphism in  $\mathcal{V}\text{-Cat}$ , certainly lies in the  $\mathcal{E}'$  of the proper factorization system  $(\mathcal{E}', \mathcal{M}')$ ; whence it follows by [6, Proposition 2.1.1] that  $F$  lies in  $\mathcal{E}'$  if and only if  $I$  lies in  $\mathcal{E}'$ . The case of a bijective-on-objects  $\mathcal{V}$ -functor, however, is dealt with in the following:

5.1. LEMMA. *A  $\mathcal{V}$ -functor  $F : \mathcal{A} \rightarrow \mathcal{B}$  which is bijective on objects lies in  $\mathcal{E}'$  if and only if it is locally in  $\mathcal{E}$ .*

PROOF. Suppose that  $F : \mathcal{A} \rightarrow \mathcal{B}$  is bijective on objects and in  $\mathcal{E}'$ ; without loss of generality we may suppose  $F$  to be the *identity* on objects. Let the  $(\mathcal{E}, \mathcal{M})$ -factorization of  $F : \mathcal{A}(A, B) \rightarrow \mathcal{B}(A, B)$  be

$$\mathcal{A}(A, B) \xrightarrow{E_{A,B}} \mathcal{D}(A, B) \xrightarrow{M_{A,B}} \mathcal{B}(A, B).$$

For objects  $A, B, C$  of  $\mathcal{A}$ , we have  $E_{B,C} \otimes E_{A,B}$  in  $\mathcal{E}$ , since the class  $\mathcal{E}$  is by assumption closed under tensor products, and we also have  $M_{A,C}$  in  $\mathcal{M}$ ; thus there is a unique map

$M'$  making commutative the diagram

$$\begin{array}{ccccc}
 \mathcal{A}(B, C) \otimes \mathcal{A}(A, B) & \xrightarrow{E_{B,C} \otimes E_{A,B}} & \mathcal{D}(B, C) \otimes \mathcal{D}(A, B) & \xrightarrow{M_{B,C} \otimes B_{A,B}} & \mathcal{B}(B, C) \otimes \mathcal{B}(A, B) \\
 \downarrow M & & \downarrow M' & & \downarrow M'' \\
 \mathcal{A}(A, C) & \xrightarrow{E_{A,C}} & \mathcal{D}(A, C) & \xrightarrow{M_{A,C}} & \mathcal{B}(A, C),
 \end{array}$$

in which  $M$  and  $M''$  are the composition maps for  $\mathcal{A}$  and  $\mathcal{B}$ . The  $M'$  give to the  $\mathcal{D}(A, B)$  the structure of a  $\mathcal{V}$ -category  $\mathcal{D}$ , for which the  $E_{A,B}$  constitute a  $\mathcal{V}$ -functor  $M : \mathcal{D} \rightarrow \mathcal{B}$  which is the identity on objects: the point is that the  $\mathcal{V}$ -category axioms for  $\mathcal{D}$  follow from those for  $\mathcal{B}$ , since the  $M_{A,B}$  are monomorphic. Now the  $E_{A,B}$  constitute a  $\mathcal{V}$ -functor  $E : \mathcal{A} \rightarrow \mathcal{D}$  which is the identity on objects, and  $F = ME$  provides a factorization of  $F$  with  $M \in \mathcal{M}'$ . Since  $F \in \mathcal{E}'$ , this implies that  $M$  is invertible, and in particular that each  $M_{A,B}$  is so, so that  $F_{A,B} = M_{A,B}E_{A,B}$  lies in  $\mathcal{E}$  as required.

Conversely, a  $\mathcal{V}$ -functor which is bijective on objects and locally in  $\mathcal{E}$  factorizes through no proper  $\mathcal{M}'$ -subobject, and so must be in  $\mathcal{E}'$ . ■

This now gives:

**5.2. PROPOSITION.** *A  $\mathcal{V}$ -functor  $F : \mathcal{A} \rightarrow \mathcal{B}$  is in  $\mathcal{E}'$  if and only if it can be written as  $F = IQ$  where  $Q$  is a quotient on objects and  $I$  is bijective on objects and locally in  $\mathcal{E}$ .*

As a first step to proving that  $\mathcal{V}\text{-Cat}$  is locally  $\lambda$ -bounded with respect to  $(\mathcal{E}', \mathcal{M}')$ , we prove:

**5.3. LEMMA.**  *$\mathcal{V}\text{-Cat}$  admits arbitrary cointersections of maps in  $\mathcal{E}'$ .*

**PROOF.** Let  $(E_i : \mathcal{A} \rightarrow \mathcal{B}_i)_{i \in I}$  be a family of  $\mathcal{V}$ -functors, each lying in  $\mathcal{E}'$ . By well-ordering the indexing set  $I$ , we can write these instead in the form  $(E_\alpha : \mathcal{A} \rightarrow \mathcal{B}_\alpha)_{\alpha < \delta}$  for some initial ordinal  $\delta$ . We set out to define by transfinite induction a “descending” family  $F_\alpha : \mathcal{A} \rightarrow \mathcal{C}_\alpha$  of  $\mathcal{V}$ -functors in  $\mathcal{E}'$ . We set  $F_0 : \mathcal{A} \rightarrow \mathcal{C}_0$  to be equal to  $1 : \mathcal{A} \rightarrow \mathcal{A}$ . We take for  $F_{\alpha+1} : \mathcal{A} \rightarrow \mathcal{C}_{\alpha+1}$  the cointersection of  $F_\alpha : \mathcal{A} \rightarrow \mathcal{C}_\alpha$  and  $E_\alpha : \mathcal{A} \rightarrow \mathcal{B}_\alpha$ ; it lies in  $\mathcal{E}'$ , because  $\mathcal{E}'$  is closed under any cointersections that exist. Finally, for a limit ordinal  $\alpha$ , we take for  $F_\alpha : \mathcal{A} \rightarrow \mathcal{C}_\alpha$  the cointersection of all the  $F_\beta$  with  $\beta < \alpha$ , provided that this exists; then  $F_\delta : \mathcal{A} \rightarrow \mathcal{C}_\delta$  is clearly the required cointersection of the  $E_\alpha : \mathcal{A} \rightarrow \mathcal{B}_\alpha$ , if it exists.

Suppose it does not. Let  $\gamma$  be the first ordinal for which  $F_\gamma$  fails to exist. Then  $\gamma$  cannot be of the form  $\alpha + 1$ , since binary cointersections certainly exist; thus  $\gamma$  is a limit ordinal. It cannot be small, since small cointersections exist. Since  $\text{ob}\mathcal{A}$  has only a small set of epimorphic images in **Set**, the surjections  $\text{ob}F_\alpha : \text{ob}\mathcal{A} \rightarrow \text{ob}\mathcal{C}_\alpha$  have become constant at some ordinal  $\beta < \gamma$ ; so that the comparison functor  $F_\rho^\sigma : \mathcal{C}_\rho \rightarrow \mathcal{C}_\sigma$  is bijective on objects whenever  $\beta \leq \rho < \sigma < \gamma$ . Since  $F_\rho^\sigma$  is in  $\mathcal{E}'$  by [6, Proposition 2.1.1], it is locally in  $\mathcal{E}$  by Lemma 5.1. But now the non-existence of the cointersection  $F_\gamma : \mathcal{A} \rightarrow \mathcal{C}_\gamma$  contradicts the hypothesis that  $\mathcal{V}_0$  admits arbitrary cointersections of maps in  $\mathcal{E}$ . ■

We have seen that  $\mathcal{V}\text{-Cat}$  is a cocomplete category with a proper factorization system  $(\mathcal{E}', \mathcal{M}')$  for which  $\mathcal{V}\text{-Cat}$  admits arbitrary  $\mathcal{E}'$ -cointersections. It will therefore be locally  $\lambda$ -bounded if it has an  $(\mathcal{E}', \mathcal{M}')$ -generator consisting of  $\lambda$ -bounded objects. Let  $\mathcal{G}$  be an  $(\mathcal{E}, \mathcal{M})$ -generator for  $\mathcal{V}_0$  consisting of  $\lambda$ -bounded objects; without loss of generality we may suppose that  $\mathcal{G}$  contains the initial object 0. Write  $\mathcal{G}'$  for the set of those  $\mathcal{V}$ -categories of the form  $\mathfrak{Z}_G$  for some  $G \in \mathcal{G}$ . The reader will easily verify that  $\mathcal{G}'$  is an  $(\mathcal{E}', \mathcal{M}')$ -generator for  $\mathcal{V}\text{-Cat}$ : the argument is essentially that used to prove Lemma 4.2. A little more work is required in showing that  $\mathfrak{Z}_G$  is  $\lambda$ -bounded in  $\mathcal{V}\text{-Cat}$  when  $G$  is so in  $\mathcal{V}_0$ , since we first need the following lemma:

5.4. LEMMA. Consider a small filtered family  $(F_j : \mathcal{A}_j \rightarrow \mathcal{B})_{j \in J}$  in  $\mathcal{M}'$ ; without loss of generality we take the functions  $\text{ob}F_j : \text{ob}\mathcal{A}_j \rightarrow \text{ob}\mathcal{B}$  to be set-inclusions. Then  $(F_j : \mathcal{A}_j \rightarrow \mathcal{B})_{j \in J}$  is an  $\mathcal{M}'$ -union in  $\mathcal{V}\text{-Cat}$  precisely when  $\text{ob}\mathcal{B}$  is the union in  $\mathbf{Set}$  of the  $\text{ob}\mathcal{A}_j$  and, for each pair  $X, Y$  of objects of  $\mathcal{B}$ , the family  $(F_j : \mathcal{A}_j(X, Y) \rightarrow \mathcal{B}(X, Y))_{j \in J_{X, Y}}$  is an  $\mathcal{M}$ -union in  $\mathcal{V}_0$ , where  $J_{X, Y}$  is  $\{j \in J \mid X \text{ and } Y \text{ lie in } \text{ob}\mathcal{A}_j\}$ .

PROOF. The  $F_j$  are an  $\mathcal{M}'$ -union if and only if the induced  $\mathcal{V}$ -functor  $F : \text{colim}_j \mathcal{A}_j \rightarrow \mathcal{B}$  lies in  $\mathcal{E}'$ . This means in particular that it is surjective on objects; but  $F$  is in any case injective on objects, since the  $F_j$  are so, and  $\mathcal{J}$  is filtered. Thus  $F$  would need to be bijective on objects, and we saw in Lemma 5.1 that such a  $\mathcal{V}$ -functor lies in  $\mathcal{E}'$  if and only if it is locally in  $\mathcal{E}$ . Thus the  $F_j$  are an  $\mathcal{M}'$ -union if and only if  $\text{ob}\mathcal{B}$  is the union of the  $\text{ob}\mathcal{A}_j$  and  $F$  is locally in  $\mathcal{E}$ .

Since  $\mathcal{J}$  is filtered, the colimit of the  $\mathcal{A}_j$  is preserved by  $U : \mathcal{V}\text{-Cat} \rightarrow \mathcal{V}\text{-Gph}$ . Thus

$$(\text{colim}_j \mathcal{A}_j)(X, Y) = \text{colim}_j \sum_{\substack{F_j A_j = X \\ F_j B_j = Y}} \mathcal{A}_j(A_j, B_j),$$

but then to say that  $F : \text{colim}_j \mathcal{A}_j(X, Y) \rightarrow \mathcal{B}(X, Y)$  is in  $\mathcal{E}$  for all  $X$  and  $Y$  is just to say that  $(F_j : \mathcal{A}_j(X, Y) \rightarrow \mathcal{B}(X, Y))_{j \in J_{X, Y}}$  is an  $\mathcal{M}$ -union in  $\mathcal{V}_0$ . ■

It now follows easily that  $\mathcal{V}\text{-Cat}(\mathfrak{Z}_G, -) : \mathcal{V}\text{-Cat} \rightarrow \mathbf{Set}$  preserves  $\lambda$ -filtered  $\mathcal{M}'$ -unions if  $\mathcal{V}_0(G, -) : \mathcal{V}_0 \rightarrow \mathbf{Set}$  preserves  $\lambda$ -filtered  $\mathcal{M}$ -unions, and so we have:

5.5. PROPOSITION.  $\mathcal{V}\text{-Cat}$  is locally  $\lambda$ -bounded with respect to  $(\mathcal{E}', \mathcal{M}')$ .

Finally, we look at the closed structure of  $\mathcal{V}\text{-Cat}$  in this context:

5.6. THEOREM. If  $\mathcal{V}$  is a symmetric monoidal closed category which is locally  $\lambda$ -bounded as a closed category with respect to the proper factorization system  $(\mathcal{E}, \mathcal{M})$ , then  $\mathcal{V}\text{-Cat}$  is locally  $\lambda$ -bounded as a closed category with respect to the proper factorization system  $(\mathcal{E}', \mathcal{M}')$ .

PROOF. We must prove for each  $\mathcal{V}$ -category  $\mathcal{X}$  that  $\mathcal{X} \otimes E$  is in  $\mathcal{E}'$  if  $E$  is so; or, equivalently, that  $[\mathcal{X}, M]$  is in  $\mathcal{M}'$  if  $M$  is so. Suppose then that  $\mathcal{X}$  is a  $\mathcal{V}$ -category and that  $M : \mathcal{A} \rightarrow \mathcal{B}$  lies in  $\mathcal{M}'$ . An object of  $[\mathcal{X}, \mathcal{A}]$  is a  $\mathcal{V}$ -functor from  $\mathcal{X}$  to  $\mathcal{A}$ ; since  $M$  is a monomorphism,  $[\mathcal{X}, M] : [\mathcal{X}, \mathcal{A}] \rightarrow [\mathcal{X}, \mathcal{B}]$  is injective on objects. To see that  $[\mathcal{X}, M]$  is locally in  $\mathcal{M}$ , let  $F, G : \mathcal{X} \rightarrow \mathcal{A}$  be  $\mathcal{V}$ -functors. Then the hom-object

$[\mathcal{X}, \mathcal{A}](F, G)$  is given by the end

$$\int_{X \in \mathcal{X}} \mathcal{A}(FA, GA).$$

Each  $M : \mathcal{A}(FX, GX) \rightarrow \mathcal{B}(MFX, MGX)$  lies in  $\mathcal{M}$ , and  $\mathcal{M}$  is closed under limits; it follows that

$$\int_{X \in \mathcal{X}} M : \int_{X \in \mathcal{X}} \mathcal{A}(FX, GX) \rightarrow \int_{X \in \mathcal{X}} \mathcal{B}(MFX, MGX)$$

lies in  $\mathcal{M}$ ; that is, that  $M : \mathcal{A}(F, G) \rightarrow \mathcal{B}(MF, MG)$  does so. ■

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